

Robustic Video Streaming in Mobile Network Using SNS Based Prefetching

¹Mrs.Meenakshi ²V.Jayashreeja ³C.Sowmya ⁴D.Gayathri

¹ASSISTANT PROFESSOR, ^{2,3,4}UG SCHOLAR, DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

¹minuraj11@gmail.com ²jayashreejavasu@gmail.com ³sowmyachinnaraj@gmail.com ⁴gayathrid33@gmail.com

^{1,2,3,4}Vel Tech High Tech Dr. Rangarajan Dr. Sakunthala Engineering College

Abstract

Due to the presence of many barriers for effective transmissions of videos in mobile networks we cannot cop up with the mobile user's requirements.

Nowadays, due to the traffic demand in the wireless communication, many problems such as poor service quality of video streaming over mobile network arise. To bare or decline such flaws, streaming technique is applied in cloud computing. On establishing various channel access links also, the gap between traffic demand and connection establishment can be bridged. The techniques such as TDMA which is based on time depending usage of the mobile users. Robustic video streaming enhances better

I. Introduction

Although proceeding our life in 21st century, the drastic development period, we experience the problem of slow video transmissions like uploading, downloading, sharing, etc.. which is due to the low bandwidth capabilities. This technology solves such problems by following two properties. While receiving video streaming traffic via 3G/4G mobile networks, mobile users often suffer from long buffering time and intermittent disruptions due to the limited bandwidth and link condition fluctuation caused by multi-path fading and user mobility. Thus, it is crucial to improve the service quality of mobile video streaming while using the networking and computing resources efficiently. Recently there have been many studies on how to improve the service quality of mobile video streaming on two aspects:

- Scalability:

This measures the video coding performance. Also, the available link capacity of a mobile device may vary over

streaming flow using scalable video coding. This technique has been advanced with SNS based prefetching in which prefetching decision can be taken based on relation between user and accessing histories recorded in client. This technique improves QoS. This SNS based prefetching enables quick buffering and avoids packet losses and intermittent disruptions during transmission. The streaming supports wide spectrum of mobile devices. This keeps track of the social network interactions under the governance of an authorized private agent. Many techniques are applied here to obtain the predicted bandwidth and proper segment transfer.

time and space depending on its signal strength, other users traffic in the same cell, and link condition variation. Storing multiple versions (with different bit rates) of the same video content may incur high overhead in terms of storage and communication. To address this issue, the Scalable Video Coding (SVC) technique (Annex G extension)

of the H.264 AVC video compression standard defines a base layer (BL) with multiple enhance layers (ELs). These substreams can be encoded by exploiting three scalability features: (i) spatial scalability by layering image resolution (screen pixels), (ii) temporal scalability by layering the frame rate, and (iii) quality scalability by layering the image compression. By the SVC, a video can be decoded/played at the lowest quality if only the BL is delivered. However, the more ELs can be delivered, the better quality of the video stream is achieved.

- Robustiveness:

Robustiveness ensures the receipt of all the layers such as base layer (BL) and enhances other layers also. Traditional video streaming techniques designed by considering relatively stable traffic links

between servers and users, perform poorly in mobile environments. Thus the fluctuating wireless link status should be properly dealt with to provide 'tolerable' video streaming services. To address this issue, we have to adjust the video bit rate adapting to the currently time-varying available link bandwidth of each mobile user. Such adaptive streaming techniques can effectively reduce packet losses and bandwidth waste. Scalable video coding and adaptive streaming techniques can be jointly combined to accomplish effectively the best possible quality of video streaming services. That is, we can dynamically adjust the number of SVC layers depending on the current link status. However most of the proposals seeking to jointly utilize the video scalability and adaptability rely on the active control on the server side. That is, every mobile user needs to individually report the transmission status (e.g., packet loss, delay and signal quality) periodically to the server, which predicts the available bandwidth for each user. Thus the problem is that the server should take over the substantial processing overhead, as the number of users increases. Cloud computing techniques are poised to flexibly provide

scalable resources to content/service providers, and process offloading to mobile users. Thus, cloud data centers can easily provision for large-scale real-time video services as investigated

in. Several studies on mobile cloud computing technologies have proposed to generate personalized intelligent agents for servicing mobile users, e.g., Cloudlet and Stratus. This is because, in the cloud, multiple agent instances (or threads) can be maintained dynamically and efficiently

depending on the time-varying user demands. Recently social network services (SNSs) have been increasingly popular. There have been proposals to improve the quality of content delivery using SNSs. In SNSs, users may share, comment or re-post videos among friends and members in the same group, which implies a user may watch a video that her friends have recommended

. Users in SNSs can also follow famous and popular users based on their interests (e.g., an official facebook or twitter account that shares the newest pop music videos), which is likely to be watched by its followers. In this regard, we are further motivated to exploit the relationship among mobile users from their SNS activities in order to prefetch in advance the beginning part of the video or even the whole video to the members of a group who have not seen the video yet. It can be done by a background job supported by the agent (of a member) in the cloud; once the user clicks to watch the video, it can instantly start playing. In this paper, we design an adaptive video streaming and prefetching framework for mobile users with the above objectives in mind, dubbed this method. Cloud constructs a private agent for each mobile user in cloud computing environments, which is used by its two main parts: (i) RMoV (robust mobile video streaming), and (ii) ESoV (efficient social video sharing). The contributions of this paper can be summarized as follows:

- RMoV offers the best possible streaming experiences by adaptively controlling the streaming bit rate depending on the fluctuation of the link quality. RMoV adjusts the bit rate for each user leveraging the scalable video coding. The private agent of a user keeps track of the feedback information on the link status. Private agents of users are dynamically initiated and optimized in the cloud computing platform.

Also the real-time SVC coding is done on the cloud computing side efficiently.

- This Cloud supports distributing video streams efficiently by facilitating a 2-tier structure: the first tier is a content delivery network, and the second tier is a data center. With this structure, video sharing can be optimized within the cloud. Unnecessary redundant downloads of popular videos can be prevented.

- Based on the analysis of the SNS activities of mobile users, ESoV seeks to provide a user with instant playing of video clips by prefetching the video clips in advance from her private agent to the local storage of her device. The strength of the social links

II. RELATED WORK

A. Robustic Video Streaming Techniques

In the adaptive streaming, the video traffic rate is adjusted on the fly so that a user can experience the maximum possible video quality based on his or her link's time-varying bandwidth capacity. There are mainly two types of adaptive streaming techniques, depending on whether the adaptivity is controlled by the client or the server. The Microsoft's Smooth

Streaming is a live adaptive streaming service which can switch among different bit rate segments encoded with configurable bit rates and video resolutions at servers, while clients dynamically request videos based on local monitoring of link quality. Adobe and Apple also developed client-side HTTP adaptive live streaming solutions operating in the similar manner. There are also some similar adaptive streaming services where servers controls the adaptive transmission of video segments, for example, the Quavlive Adaptive Streaming. However, most of these solutions maintain multiple copies of the video content with different bit rates, which brings huge burden of storage on the server. Regarding rate adaptation controlling techniques, TCP friendly rate control methods for streaming services over mobile networks are proposed, where TCP throughput of a flow is predicted as a function of packet loss rate, round trip time, and packet size. Considering the estimated throughput, the bit rate of the streaming traffic can be adjusted. A rate adaptation algorithm for conversational 3G video streaming is introduced by. Then, a few

between users and the history of various social activities can probabilistically determine how much and which video will be prefetched. The rest of the paper is organized as follows. We first introduce related work in Section II, and explain the Cloud framework in Section III. The adaptive video streaming service and the efficient social video sharing will be detailed in Sections IV and V, respectively. Then the operations of this Cloud is illustrated in Section VI. Finally, we evaluate the prototype implementation in Section VII, and conclude the paper in Section VIII.

cross-layer adaptation techniques are discussed, which can acquire more accurate information of link quality so that the rate adaptation can be more accurately made. However, the servers have to always control and thus suffer from large workload. Recently the H.264 Scalable Video Coding (SVC) technique has gained a momentum. An adaptive video streaming system based on SVC is deployed in, which studies the real-time SVC decoding and encoding at PC servers. The work

in proposes a quality-oriented scalable video delivery

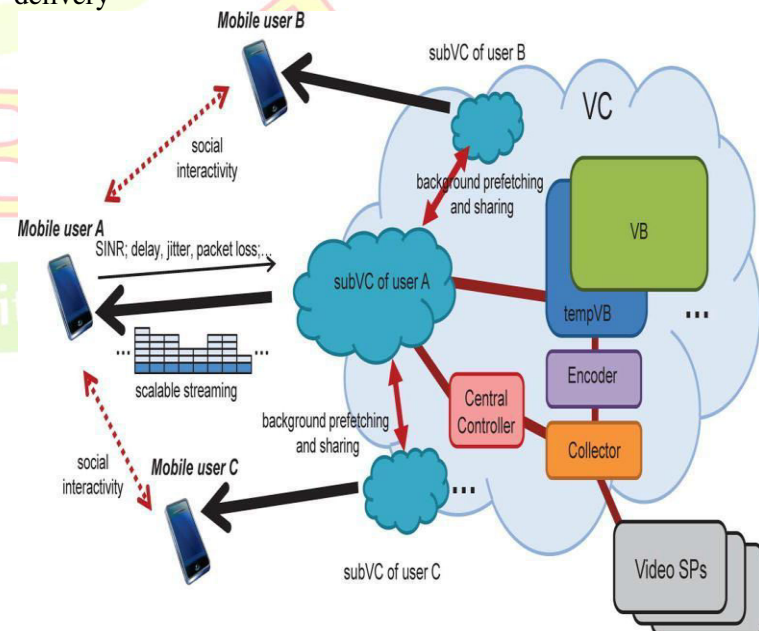


Fig. 1. Illustration of the AMES-Cloud framework with the Video Cloud (VC), subVCs

for mobile users, the Video Base (VB), and the Video Service Providers (SPs).

using SVC, but it is only tested in a simulated LTE Network. Regarding the encoding performance of SVC, CloudStream mainly proposes to deliver high-quality streaming videos through a cloud-based SVC proxy, which discovered that the cloud computing can significantly improve the performance

of SVC coding. The above studies motivate us to use SVC for video streaming on top of cloud computing.

B. Mobile Cloud Computing Techniques

The cloud computing has been well positioned to provide video streaming services, especially in the wired Internet because of its scalability and capability. For example, the quality-assured

III. CLOUD FRAMEWORK

In this section we explain the AMES-Cloud framework includes the Adaptive Mobile Video streaming (AMoV) and the Efficient Social Video sharing (ESoV). As shown in Fig. 1, the whole video storing and streaming system in the cloud is called the Video Cloud (VC). In the VC, there is a large-scale video base (VB), which stores the most of the popular video clips for the video service providers (VSPs). A temporal video base (tempVB) is used to cache new candidates for the popular videos, while tempVB counts the access frequency of each video. The VC keeps running a collector to seek videos which are already popular in VSPs, and will re-encode the collected videos into SVC format and store into tempVB first. By this 2-tier storage, the AMES-Cloud can keep serving most of popular videos eternally. Note that management work will be handled by the controller in the VC. Specialized for each mobile user, a sub-video cloud (subVC) is created dynamically if there is any video streaming demand from the user. The sub-VC has a sub video base (subVB), which stores the recently fetched video segments. Note that the video deliveries among the subVCs

IV. AMOV: ADAPTIVE MOBILE VIDEO STREAMING

A. SVC

As shown in Fig. 2, traditional video streams with fixed bitrates cannot adapt to the fluctuation of the link quality. For a particular bitrate, if the sustainable link bandwidth

bandwidth auto scaling for VoD streaming based on the cloud computing is proposed, and the CALMS framework is a cloud-assisted live media streaming service for globally distributed users. However, extending the cloud computing-based services to mobile environments requires more factors to consider: wireless link dynamics, user mobility, the limited capability of mobile devices. More recently, new designs for users on top of mobile cloud

computing environments are proposed, which virtualize private agents that are in charge of satisfying the requirements (e.g., QoS) of individual users such as Cloudlets and Stratus. Thus,

we are motivated to design the Cloud framework by using virtual agents in the cloud to provide adaptive video streaming services.

and the VC in most cases are actually not —copyl, but just —linkl operations on the same file eternally within the cloud data center. There is also encoding function in subVC (actually a smaller-scale encoder instance of the encoder in VC), and if the mobile user demands

a new video, which is not in the subVB or the VB in VC, the subVC will fetch, encode and transfer the video. During video streaming, mobile users will always report link conditions to their corresponding subVCs, and then the subVCs offer adaptive video streams. Note that each mobile device also has a temporary caching storage, which is called local video base (localVB), and is used for buffering and prefetching. Note that as the cloud service may access different places, or even continents, so in the case of a video delivery and prefetching between different data centers, a transmission will be carried out, which can be then called —copyl. And because of the optimal deployment of data centers, as well as the capable

links among the data centers, the —copyl of a large video file takes tiny delay.

varies much, the video streaming can be frequently terminated due to the packet loss. In SVC, a combination of the three lowest scalability is called the Base Layer (BL) while the enhanced combinations are called

Enhancement Layers (ELs). To this regard, if BL

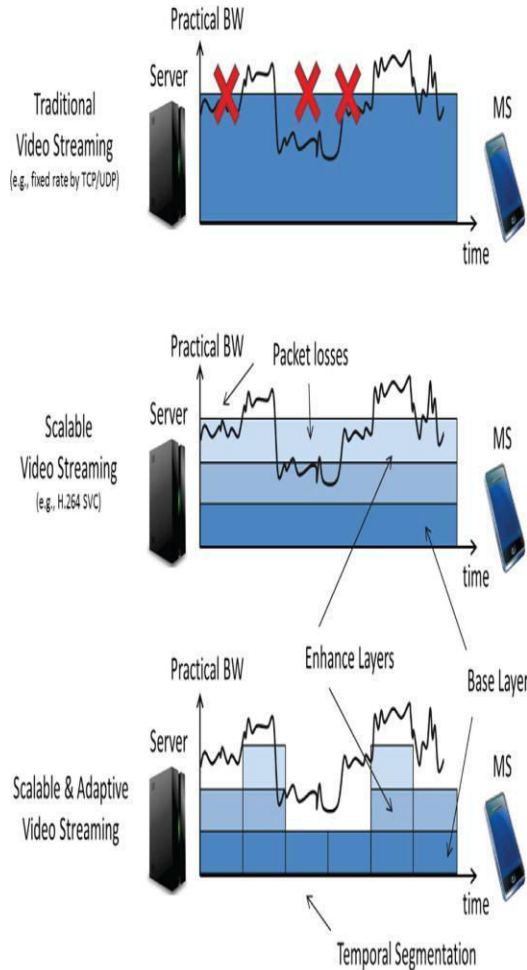


Fig. 2. A comparison of the traditional video streaming, the scalable video streaming and the streaming in the AMES-Cloud framework.

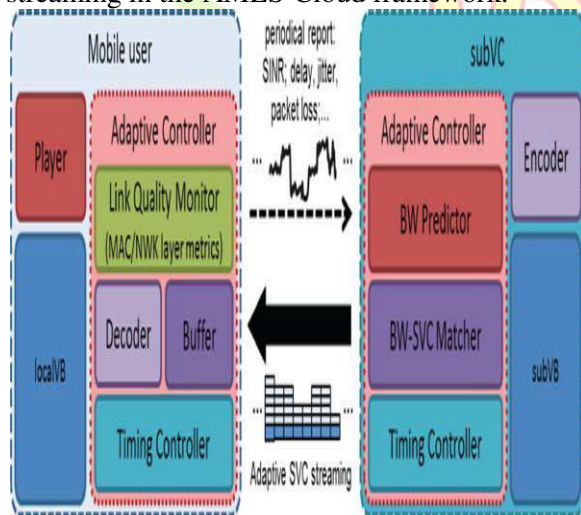


Fig. 3. Functional structure of the client and the subVC.

is guaranteed to be delivered, while more ELs can be also obtained when the link can afford, a better video quality can be expected. By using SVC encoding techniques, the server doesn't need to concern the client side or the link quality. Even some packets are lost, the client still can decode the video and display. But this is still not bandwidth-efficient due to the unnecessary packet loss. So it is necessary to control the SVC-based video streaming at the server side with the rate adaptation method to efficiently utilize the bandwidth.

B. Adaptability With Monitoring on Link Quality

We design the mobile client and the subVC with the structure as shown in Fig. 3. The link quality monitor at mobile client keeps tracking on metrics including signal strength, packet round-trip-time (RTT), jitter and packet loss with a certain duty cycle. And the client will periodically report to the subVC.

Hereby we define the cycle period for the reporting as the Δt , denoted by Δt .

Note that the video is also split by temporal segmentation by interval. Once the subVC gets the information of the link quality, it will perform a calculation and predict the potential bandwidth in the next time window.

Note that we will use \hat{B}_t —predicted bandwidth and \hat{G}_t —predicted goodput interchangeably in following parts.

Suppose sequence number of current time window is n , the predicted bandwidth can be estimated by:

where, α indicating the importance of each factor, β is for packet loss rate, γ is for RTT, δ is for signal to interference and noise ratio, and ϵ , η , θ are three functions reflecting the value change of each factor compared

with that of last time window. Actually in this paper we deploy a measurement-based prediction,

that is we directly use of last time window as the of next time window, which is proved with already high accuracy.

C. Matching Between Bandwidth Prediction and SVC Segments

After obtaining the predicted bandwidth, or say goodput, of next time window, subVC will match and decide how many video segments of BL and ELs can be transmitted approximately. We hereby define the term —resolutionl to indicate the level of temporal segmentation and the number of ELs. If is small and there are more ELs, we say the SVC-based video source is with a higher resolution. We illustrate two cases of

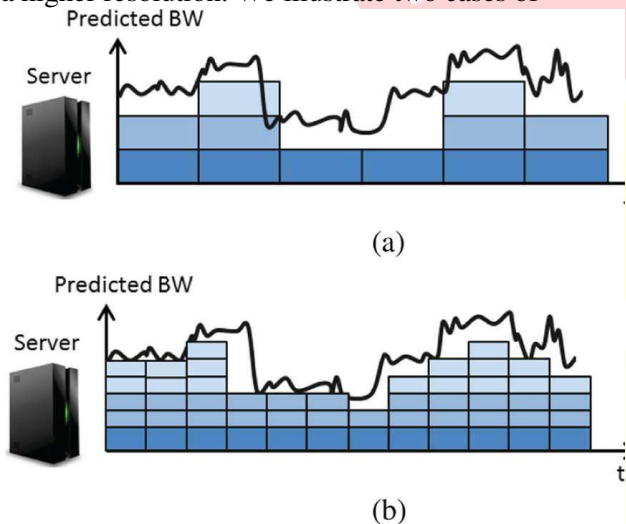


Fig. 4. Matching between predicted bandwidth and SVC-segments with different resolutions. (a) Fine-grained (high resolution). (b) Coarse-grained (low resolution).

coarse-grained (low resolution) and a relatively fine-grained (high resolution) for matching between the SVC segments and the predicted goodput in Fig. 4. The resolution with two ELs and a larger can hardly fit to the signal fluctuation, and thus there are some bandwidth wasted or packets lost. In contrast a higher

V. ESOV: EFFICIENT SOCIAL VIDEO SHARING

A. Social Content Sharing

In SNSs, users subscribe to known friends, famous people, and particular interested content publishers as well; also there are various types of social activities among users in SNSs, such as direct message and public posting. For spreading videos in SNSs, one can post a video in the public, and his/her subscribers

can quickly see it; one can also directly recommend a video to specified friend(s); furthermore one can periodically get noticed by subscribed content publisher for new or popular

resolution with more ELs and a smaller can always fit the fluctuation of the bandwidth. However a higher resolution also induces more encoding workload to the servers. Suppose there are totally ELs, and the bit rate of the th EL is denoted as while the bit rate of the BL is). We let indicate the SVC segment of BL with temporal sequence, and let indicate the SVC segment of the th EL with temporal sequence . So the algorithm of matching between predicted bandwidth and SVC segments is shown in Algorithm 1 as following:

Algorithm 1 Matching Algorithm between BW and Segments

```

i=0
BW0=RBL
Transmit BL0
Monitor BW0practical
repeat
Sleep for Twin
Obtain ,pi,RTTi,SINRi etc., from client's report
Predict BWi+1estimate (or BWi+1estimate=BWipractical)
k=0
BWEL=0
repeat
k++
if k>=j break
BWEL=BWEL+RELk
until BWEL>=BWi+1estimate-RBL2
Transmit BLi+1 and ELi+11,ELi+12,...,ELi+1k-1
Monitor BWi+1practical
i++
until All video segments are transmitted

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videos. Similar to studies in, we define different strength levels for those social activities to indicate the probability that the video shared by one user may be watched by the receivers of the one's sharing activities, which is called a

—hitting probabilityl, so that subVCs can carry out effective background prefetching at subVB and even localVB. Because after a video sharing activity, there may be a certain delay that the recipient gets to know the sharing, and initiates to watch. Therefore the prefetching in prior will not impact the users at most cases. Instead, a user

can click to see without any buffering delay as the beginning part or even the whole video is already prefetched at the local VB. The amount of prefetched segments is mainly

determined by the strength of the social activities. And the prefetching from VC to subVC only refers to the —linking action, so there is only file locating and linking operations with tiny delays; the prefetching from subVC to local VB also depends on the strength of the social activities, but will also consider the wireless link status. We classify the social activities in current popular SNSs into three kinds, regarding the impact of the activities and the potential reacting priority from the point of view of the recipient:

VI. Influence of SNS in the domain of mobile users:

People use social networking sites for meeting new friends, finding old friends, or locating people who have the same problems or interests they have, called niche networking. More and more relationships and friendships are being formed online and then carried to an offline setting. Psychologist and University of Hamburg professor Erich H. Witte says that relationships which start online are much more likely to succeed. Witte has said that in less than 10 years, online dating will be the predominant way for people to start a relationship. One online dating site claims that 2% of all marriages begin at its site, the equivalent of 236 marriages a day. Other sites claim one in five relationships begin online. Users do not necessarily share with others the content which is of most interest to them, but rather that which projects a good impression of themselves. While everyone agrees that social networking has had a significant impact on social interaction, there remains a substantial disagreement as to whether the nature of this impact is completely positive. A number of scholars have done research on the negative effects of internet communication as well. These researchers have contended that this form of communication is an impoverished version of conventional face-to-face social interactions, and therefore produce negative outcomes such as loneliness and depression for users who rely on social networking entirely. By engaging solely in online communication, interactions between

communities, families, and other social groups are weakened.

B. Prefetching Levels

Different strengths of the social activities indicate different levels of probability that a video will be soon watched by the recipient. Correspondingly we also define three prefetching levels

regarding the social activities of mobile users:

- —Partsl: Because the videos that published by subscriptions may be watched by the subscribers with a not high probability, we propose to only push a part of BL and EL segments, for example, the first 10% segments.

- —Alll: The video shared by the direct recommendations will be watched with a high probability, so we propose to prefetch the BL and all ELs, in order to let the recipient(s) directly watch the video with a good quality, without any buffering.

- —Littlel: The public sharing has a weak connectivity among users, so the probability that a user's friends

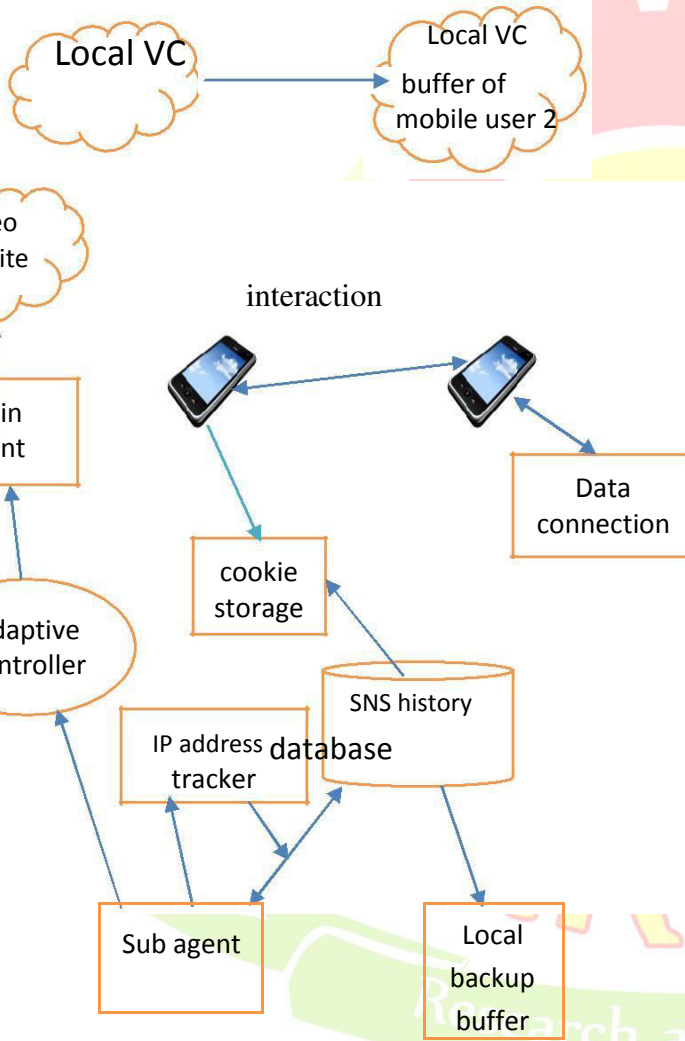
(followers) watch the video that the user has watched or shared is low. We propose to only prefetch the BL segment of the first time window in the beginning to those who have seen his/her activity in the stream. The prefetching happens among subVBs and the VB,

also more importantly, will be performed from the subVB to local VB of the mobile device depending on the link quality. If a mobile user is covered by Wi-Fi access, due to Wi-Fi's capable link and low price (or mostly for free), subVC can push as much as possible in most cases. However if it is with a 3G/4G connection, which charges a lot and suffers limited bandwidth, we propose to

downgrade the prefetching level to save energy and cost as listed in Table I, but users can still benefit from the prefetching effectively. Note that some energy prediction methods can be deployed in order to actively decide whether current battery status is suitable for —partsl or —littlel. If a user, A, gets the direct recommendation of a video from another user, B,

A's subVC will immediately prefetch the video either from B's subVB, or from the VB (or tempVB) at the level of —alll, if A is with Wi-Fi access. However if user A is connected to 3G/4G link, we will selectively prefetch a part of the

video segment to A's local storage at the level of parts. Note that the subscribed videos will be not prefetched when user A is at 3G/4G connection, as it is downgraded from little to none. A better extension of the prefetching strategy by social activities can be designed by PREFETCH MECHANISM:



an self-updating mechanism from the user's hitting history in an evolutionary manner. This learning based prefetching is out of the scope of this paper, and will be explored as our future work.

relation to placing too much personal information in the hands of large corporations or governmental bodies, allowing a profile to be produced on an individual's behavior on which decisions, detrimental to an individual, may be taken. Furthermore, there is an issue over the control of data—information that was altered or removed by

the user may in fact be retained and passed to third parties. This danger was highlighted when the controversial social networking site Quechup harvested e-mail addresses from users' e-mail accounts for use in a spamming operation. In medical and scientific research, asking subjects for information about their behaviors is normally strictly scrutinized by institutional review boards, for example, to ensure that adolescents and their parents have informed consent. It is not clear whether the same rules apply to researchers who collect data from social networking sites. These sites often contain a great deal of data that is hard to obtain via traditional means. Even though the data are public, republishing it in a research paper might be

considered invasion of privacy. Privacy on social networking sites can be undermined by many factors. For example, users may disclose personal information, sites may not take adequate steps to protect user privacy, and third parties frequently use information posted on social networks for a variety of purposes. "For the Net generation, social networking sites have become the preferred forum for social interactions, from posturing and role

playing to simply sounding off. However, because such forums are relatively easy to access, posted content can be reviewed by anyone with an interest in the users' personal information". Following plans by the UK government to monitor traffic on social networks schemes similar to e-mail jamming have been proposed for networks such as Twitter and Facebook. These would involve "friending" and "following" large numbers of random people to thwart attempts at network analysis. Privacy concerns have been found to differ between users according to gender and personality. Women are less likely to publish information that reveals methods of

VII. Issues

A. Privacy

Privacy concerns with social networking services have been raised growing concerns amongst users on the dangers of giving out too much personal information and the threat of sexual predators. Users of these services also need to be aware of data theft or viruses. However, large services, such as MySpace and Netlog, often work with law enforcement to try to prevent such incidents. In addition, there is a perceived privacy threat in

contacting them. Personality measures openness, extraversion, and conscientiousness were found to positively affect the willingness to disclose data,

VIII. VIDEO STORAGE AND STREAMING FLOWBY AMOV AND EMOS

The two parts, AMoV and EMoS, in AMES-Cloud framework have tight connections and will together service the videostreaming and sharing: they both rely on the cloud computing platform and are carried out by the private agencies of users; while prefetching in EMoS, the AMoV will still monitor and improve the transmission considering the link status; with a certain amount of prefetched segments by EMoS, AMoV can offer better video quality. With the efforts of AMoV and EMoS, we illustrate the flowchart of how a video will be streamed in Fig. 5. Note that in order to exchange the videos among the local VBs, subVBs, tempVB and the VB, a video map (VMap) is used to indicate the required segments. Once a mobile user starts to watch a video by a link, the local VB will first be checked whether there is any prefetched segments of the video so that it can directly start.

IX. IMPLEMENTATION AND EVALUATION

We evaluate the performance of the AMES-Cloud framework by a prototype implementation. We choose the U-cloud server (premium) in the cloud computing service offered by Korean Telecom, and utilize the virtual server with 6 virtual CPU cores (2.66 GHz) and 32 GB memory, which is fast enough for encoding 480P (480 by 720) video with H.264 SVC format in 30 fps at real time [9]. In the cloud, we deploy our server application based on Java, including one main program handling all tasks of the whole VC, while the program dynamically initializes, maintains and terminates instances of another small Java application as private agents for all active users. We implement the mobile client at a mobile phone, Samsung Galaxy II, with android system version 4.0. The mobile data service is offered by LG LTE network, while in some uncovered area the 3G network is used.

Note that we still use —3G| to indicate the general cellular network. We test in the downtown area, so the practical bandwidth of the mobile

while neuroticism decreases the willingness to disclose personal information.

If there is none or just some parts, the client will report a corresponding V Map to its subVC. If the subVC has prefetched parts in subVB, the subVC will initiate the segment transmission. But if there is also none in the subVB, the tempVB and VB in the center VC will be checked. For a non-existing video in AMES-Cloud, the collector in VC will immediately fetch it from external video providers via the link; after re-encoding the video into SVC format, taking a bit longer delay, the subVC will transfer to the mobile user. Also in AMES-Cloud, if a video is shared among the subVCs at a certain frequency threshold (e.g., 10 times per day), it will be uploaded to the tempVB of the VC; and if it is further shared at a much higher frequency (e.g., 100 times per day), it will be stored with a longer lifetime in the VB. In such a manner, which is quite similar to the leveled CPU cache, the subVB and VB can always store fresh and popular videos in order to increase the probability of re-usage.

link is not as high as we expected, but this won't impact our experiment results. The test video is the Tomb Raider 2012 Trailer in H.264 format with 480P resolution downloaded from YouTube. Its size

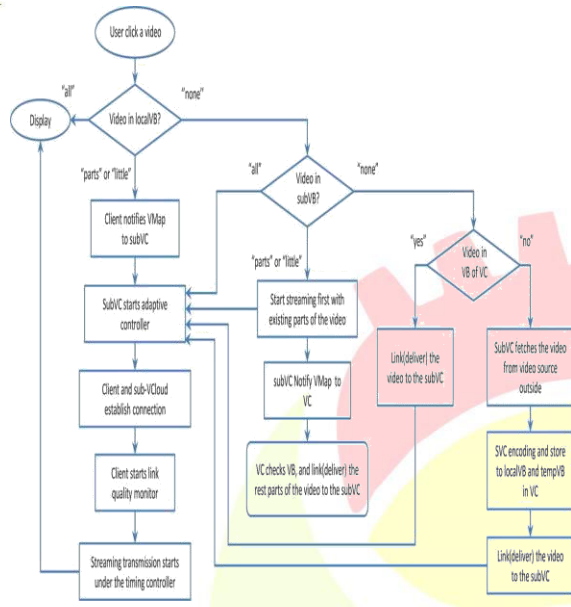


Fig. 5. Working flow of video streaming in the subVC and VC of AMES-Cloud framework.

is 13.849 Mbytes and with a duration of 180 seconds. We first decode it by the x264 decoder into the YUV format, and re-encode it by the H.264 SVC encoder, the Joint Scalable Video Model (JSVM) software of version 9.1. We just use default settings for the decoding and encoding, and do the H.264

SVC encoding at the virtual server in the cloud. We split the video into segments by 1 second to 5 seconds, that is to vary T_{win} with values 1s, 2s, 3s, 4s and 5s. By JSVM, besides the base layer, we further make five temporal layers (1.875, 3.75, 7.5, 15, and 15 fps), two spatial layers (240 by 360 and 120 by 180) and two more quality layer (low and high), referring to [12] and [40]. Thus we define the best resolution configuration as $-1+5+2+2$. And we also test different resolution configurations, including $-1+1+1+1$, $-1+2+2+2$, $-1+3+2+2$ and $-1+4+2+2$.

A. Adaptive Video Streaming Based on SVC

Firstly we examine whether there is a deep relationship between the measured bandwidth of last time window and the practical bandwidth of next time window (goodput by Kbps). We test the video streaming service via cellular link, and move the device around in the building to try to change the signal quality. Note that all tests are ran five times. The collected the relative errors

for the predicted bandwidth to the practical bandwidth for every time window, calculated by $|BW_{estimate} - BW_{practical}| / \{BW_{practical}\}$, are shown in Fig. 6, where the bar indicates the 25% and 75% quartiles, and the whiskers indicate the 5% and 95% percentiles. When T_{win} is 1 second or 2 seconds, the predicted bandwidth is very near to the practical one with around 10% relative error, but large values of T_{win} have relatively poor prediction accuracy,

which reflects the similar results. So we suggest a short T_{win} of 2 or 3 seconds for accurate prediction in practical designs

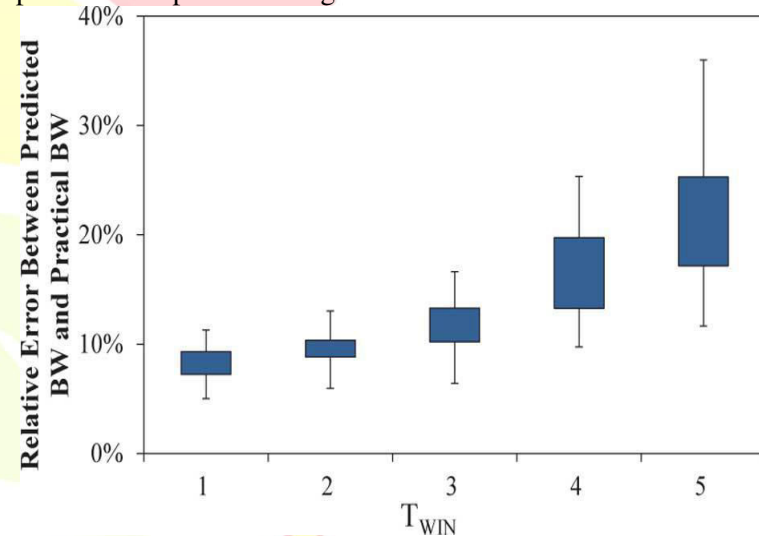


Fig. 6. Relative errors between predicted bandwidth and practical bandwidth (percentage).

B. Video Streaming in SubVC and VC

We evaluate how H.264 SVC works in AMES-Cloud framework regarding the above mentioned SVC resolution configurations. As shown in Fig. 7(a), because of the strong computational capacity by the cloud computing, the encoding speed is fast. The best resolution configuration $-1+5+2+2$ with 5

second temporal segmentation scheme requires about 560 ms for encoding. For shorter intervals of T_{win} , the encoding delay is very small under 50 ms. Because more ELs induce higher overhead due to the duplicated I-frames, we test the overhead, which is calculated by the ratio of the total size of the video segments after SVC encoding to the size of only the BL. As shown in Fig. 7(b), the resolution scheme of

$-1+1+1+1$ has a low overhead around below 10%, and $-1+2+2+2$ with two ELs for

each scalability feature has about 17% overhead, which is acceptable. However higher resolution like $-1+4+2+2$ has 61% overhead, and $-1+5+2+2$ has even 120% overhead, which is not efficient. Overall, an SVC stream should not contain too many enhance layers for extremely high scalability, which may practically bring too much overhead.

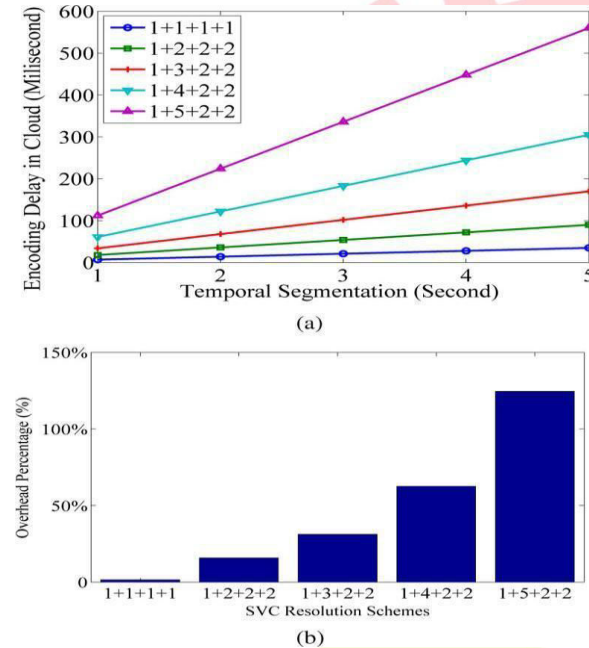


Fig. 7. Evaluation of SVC resolution schemes. (a) Delay of difference SVC resolution schemes in the Cloud (b) Overhead of different SVC resolution schemes in the Cloud.

TABLE II
DELAYS OF PREFETCHING SHARING FOR VARIOUS LEVELS

	Little	Parts
subVBs ↔ VB	0.011 s	0.023 s
subVB → locVB via Wi-Fi	2.421 s	4.359 s
subVB → locVB via 3G	N&A	18.430 s (little)

C. Prefetching Delays

In ESoV, video segments can be prefetched among VB, tempVB, and localVBs of the mobile users, based on their activities in SNSs. we

evaluate the required delays for different levels of prefetching as shown in Table. 2. We here use the normal resolution configuration of $-1+2+2+2$ with 2second temporal segmentation by default (the same in following tests). We also set the sharing length of $-little$ as only the first 5 seconds of the BL and ELs, that of $-parts$ as the first 15 seconds of the BL and ELs, and that of $-all$ as all BL and ELs segments. We can see that prefetching supported by the cloud computing is significantly fast. When prefetching via wireless links, it takes several seconds. However it is obvious that in most cases, a recipient of the video sharing may not watch immediately after the original sharing behavior, that is normal users have significant access delay gaps, so this prefetching transmission delay won't impact user's experience at all, but will bring $-non-buffering$ experience in fact when the user clicks to watch at a later time.

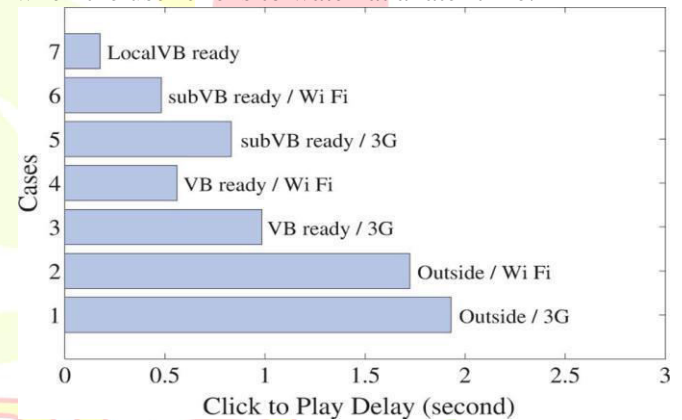


Fig. 8. Average click-to-play delay for various cases.

D. Watching Delay

We test how long one user has to wait from the moment that one clicks the video in the mobile device to the moment that the first streaming segment arrives, which is called as $-click-to-play$ delay. As shown in Fig. 8, if the video has been cached in localVB, the video can be displayed nearly immediately with ignorable delay. When we watch video which is fetched from the subVC or the VC, it generally takes no more than 1 second to start. However if the user accesses to AMES-Cloud service via the cellular link, he will still suffer a bit longer delay (around 1s) due to the larger RTT of transmission via the cellular link.

For the cases to fetch videos which are not in the AMES Cloud (but in our server at lab), the delay

is a bit higher. This is mainly due to the fetching delay via the link from our server at lab to the cloud data center, as well as the encoding delay. In practice, there are optimized links in the Internet backbone among video providers and cloud providers, and even recent video providers

are just using cloud storage and computing service. Therefore this delay can be significantly reduced in practice. Also this won't happen frequently, since most of the popular videos will be already prepared in the AMES-Cloud.

Conclusion

In this paper, we discussed our proposal of an robust mobile video streaming and social network interactions, which is our advanced robust video streaming efficiently stores videos in the clouds (VC), and utilizes cloud computing to construct private agent (subVC) for each mobile user to try to offer —ALBV (Anytime Low-Buffered Video). With this technique the user's communication in the virtual world is identified and predicted to fulfill their downloading or online procedures with the storage buffers. Privacy concerns with social networking services have been raised growing concerns amongst users on the dangers of giving out too much personal information and the threat of sexual predators. Users of these services also need to be aware of data theft or viruses. However, large services, such as MySpace and Netlog, often work with law enforcement to try to prevent such incidents. We evaluated the prefetching by prototype implementation and shows that the cloud computing technique with the involvement of SNS brings significant improvement on the robustness of the mobile streaming. The objective is to verify how cloud computing can improve the video transmission and low latency for mobile users. In the future, we will also try to improve the security issues in the Cloud.

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